Nature of numerical computations

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Abstract

Contents of the lecture:

- ☞ Absolute and relative errors.
- ☞ Number representation.

Absolute and relative errors

Let *x*0 be some real number than we need to compute. Let *x* denotes the result of computations.

 $-$ Typeset by FoilTEX $-$

Relative error: example

Consider the expression

$$
\frac{\left|e^x - 1 - x - \frac{x^2}{2!} - \dots - \frac{x^n}{n!}\right|}{|e^x|},
$$

i.e., the relative error of the Taylor approximation. The theory says that for every *x* this expression goes to 0 an *n* increases.

Consider the next script:

```
nTerms = 50;for x = [10 \ 5 \ 1 \ -1 \ -5 \ -10]figure;
    term = 1; s = 1;
    f = exp(x) * ones(nTerms, 1);for k = 1:nTermsterm = x.*term/k;s = s + term;err(k)=abs(f(k)-s);end
    relerr = err/exp(x);
    semilogy(1:nTerms, relerr);
    ylabel('Relative error in partial sum');
    xlabel('Order of partial sum');
    title(sprintf('x = %5.2f', x));
end
```
 $-$ Typeset by FoilT E^X – 1

Figure 1: Error in Taylor approximations to e^x , $x = -1$

The script proves that the relative error does not go to 0 as the number of terms in the series increases. How to explain this phenomenon?

Number presentation

Usually we represent numbers on a decimal base, for example

$$
1957 = 1 \times 10^3 + 9 \times 10^2 + 5 \times 10^1 + 7 \times 10^0
$$

 $-$ Typeset by FoilT_EX $-$ 2

The fractional part is presented by using negative powers, for example

$$
0.14 = 1 \times 10^{-1} + 4 \times 10^{-2}.
$$

The elements of hardware can not have 10 possible positions, but only 2 positions. On a computer we must use binary base, for example

 $0.5 = 2^{-1}$

Since only a finite memory space is available to store the number, we will have a rounding error.

Number presentation in memory

Sixty four bits are used to represent a real number in MATLAB:

☞ 52 bits are used to represent the fractional part, f;

☞ 11 bits are used to represent the biased exponent, e;

☞ 1 bit is used to represent the sign, s.

Figure 2: MATLAB real number format

The real number *x* presents as the machine number

$$
x = \pm 2^p \sum_{k=1}^{52} \alpha_k \cdot 2^{-k},\tag{1}
$$

where every α_k is equal to 0 or 1 and lies in the fractional part. The number p is presented according to the next rules: according to the next rules:

$$
- \text{Typeset by Foi} \Pi_{\text{E}} \text{X} - \text{3}
$$

Table 1: Presentation of exponent

If x is the machine number (1), then the next larger machine number is

$$
x_{+} = \pm 2^{p} \left(\sum_{k=1}^{52} \alpha_{k} \cdot 2^{-k} + 2^{-52} \right).
$$

The spacing of these numbers, or the absolute rounding error, is

$$
x_{+} - x = 2^{-52} \cdot 2^{p} = 2^{p-52}.
$$

The relative spacing, or the relative rounding error, is defined by the ratio

$$
\frac{x_+ - x}{|x|} \approx 2^{-52} \approx 2.2204 \times 10^{-16}.
$$

If $x = 1$, then the absolute rounding error is equal to the relative rounding error, and both of them are equal to the distance from 1 to the next machine number. This is what we call eps in MATLAB.

Conclusions

- ☞ The set of numbers in a computer is finite, since rounding errors.
- ☞ The spacing of computer numbers is not uniform, but the relative roundoff error is constant.
- ☞ Algorithms that are equivalent mathematically may behave very different numerically.

Problems

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- 1. The approximation for the cosine function is given by

$$
T_n(x) = 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \dots + (-1)^n \frac{x^{2n}}{(2n)!}.
$$

Write M-file that explores the relative error in the partial sums. Use the program at p. 1 as a pattern.

2. The next two functions

$$
y_1(x) = \sqrt{x^2 + 1} - 1
$$
 and $y_2(x) = \frac{x^2}{\sqrt{x^2 + 1} + 1}$

are mathematically equivalent. Plot the function $|y_1(x) - y_2(x)|$ over increasingly smaller neighbourhoods around *x* = 0. Use x=linspace(-delta,delta,n) for *n* = 100 and several values of delta in the interval $[10^{-4}, 10^{-3}]$.

3. (For pass with distinction). Explain the results obtained in Problem 2.